

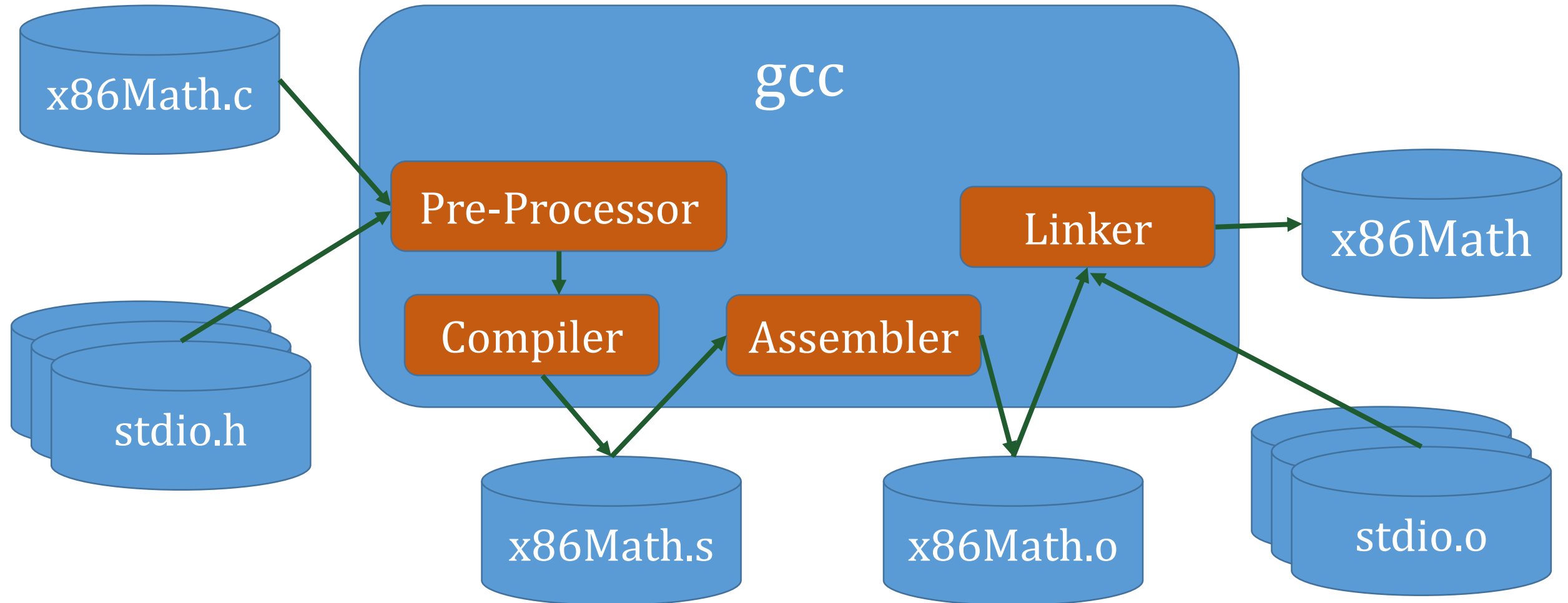
Loading Code

Computer Systems Chapter 7.5, 7.8, 7.9

```
gcc -g -o ttt ttt.c
```



```
gcc -g -o x86Math x86Math.c
```



What is in a binary executable file?

- Binary representation of X86 instructions
 - “objdump -d x86Math” disassembles these and writes them out

00000000000006f0 <main>:

6f0: 55
6f1: 48 89 e5
6f4: 48 83 ec 20
6f8: 89 7d ec
6fb: 48 89 75 e0
6ff: 83 7d ec 01
703: 7f 25
...

push %rbp
mov %rsp,%rbp
sub \$0x20,%rsp
mov %edi,-0x14(%rbp)
mov %rsi,-0x20(%rbp)
cmpl \$0x1,-0x14(%rbp)
jg 72a <main+0x3a>

Binary X86

Disassembly

- What else is in the binary executable file? ...

What else is in a binary file?

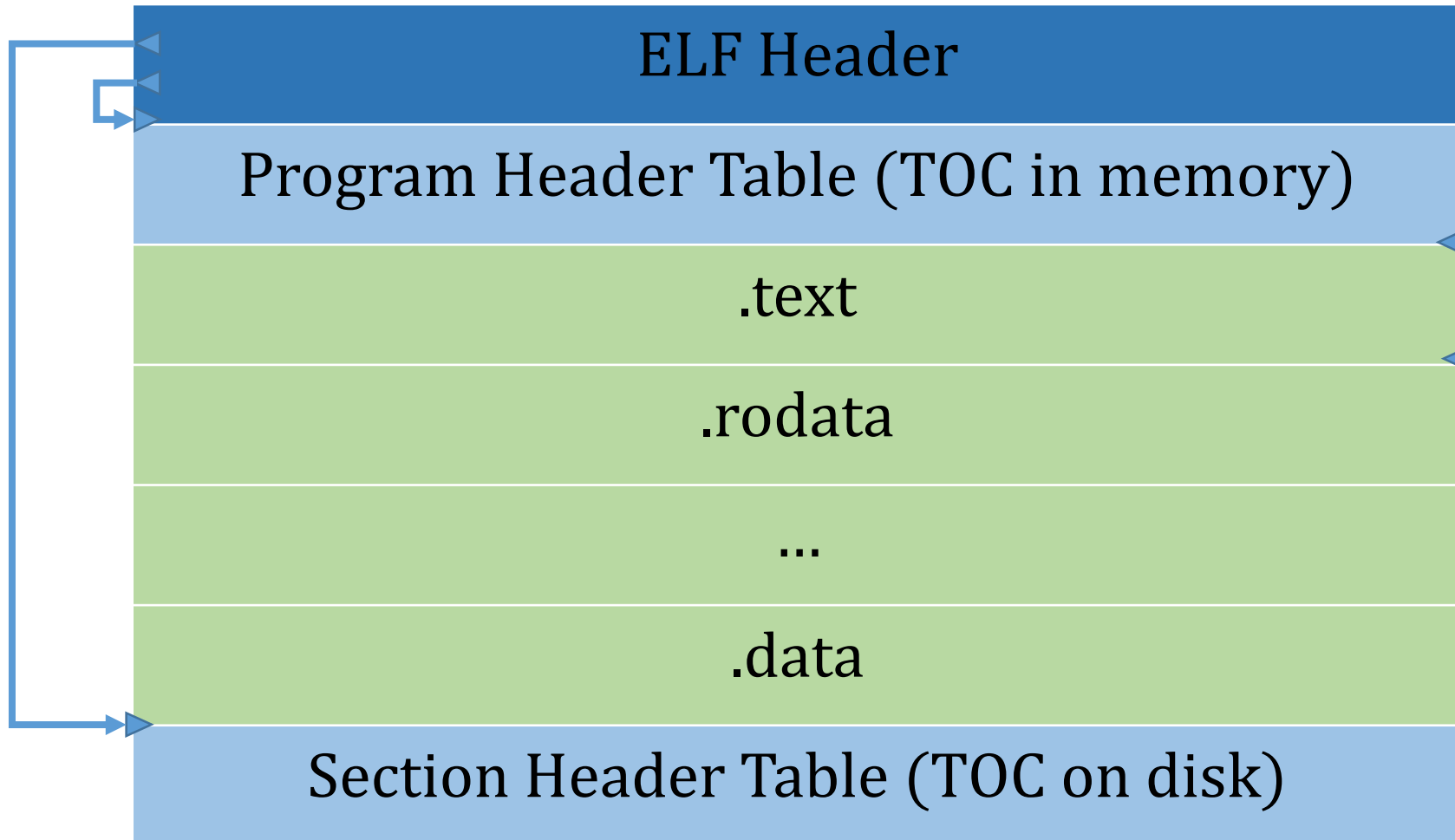
- Information about WHERE in memory the code is placed
- Cross-reference between function name and location in memory
- Information about constants
 - Some constants can be literal values in X86 instructions... \$12
 - Not all constants fit in instructions... "Invoke as %s <number>\n"
 - Binary file must contain both the value and location of constants
- Information about global variables
 - Where each global variable exists in memory
 - What the initial value of the global variable is (if initialized)
- All the debug information created by -g

Object/Executable Code ELF format

- ELF –Acronym from: Executable and Linkable Format
- First defined in 1983 UNIX “System V”
- Used for many different architectures (very popular)
- In 1996, chosen as standard for X86
- See

https://en.wikipedia.org/wiki/Executable_and_Linkable_Format

ELF File Sections



ELF Header

- “Magic Number” – First 4 bytes identify this as ELF (0x7f + ‘ELF’)

x7f	x45	x4C	x46
.	E	L	F

- Information about this file:
 - 32/64 bit addresses, big/little endian
 - Target operating system and architecture
 - relocatable, executable, shared, or core
- Program Table Info (loc, size, #entries) for use after load
- Offset of first instruction
- Section Table Info (offset, size, #entries) File table of contents

“Reading” ELF files : objdump

- -f : Interpret ELF header
- -h : List section headers (table of contents)
- -d : Disassemble x86 binary code (.text) segment
- -s : dump everything in hex
- -j<section> to restrict to a specific section
- -t : print symbol table

ELF Header Information

```
> objdump -f x86Math
```

```
x86Math:    file format elf64-x86-64
```

```
architecture: i386:x86-64, flags 0x00000150:
```

```
HAS_SYMS, DYNAMIC, D_PAGED
```

```
start address 0x000000000000005c0
```

ELF file Sections

- Need different sections for different types of data
- Each section has it's own internal data format
- ELF header points to section table
- Section Table keeps “Section Header” for each segment

objdump -h x86Math (section headers)

Idx	Name	Size	VMA	LMA	File off	Algn
0	.interp	0000001c	00000000000000238	00000000000000238	00000238	2**0
	CONTENTS, ALLOC, LOAD, READONLY, DATA					
...						
10	.init	00000017	00000000000000568	00000000000000568	00000568	2**2
	CONTENTS, ALLOC, LOAD, READONLY, CODE					
11	.plt	00000030	00000000000000580	00000000000000580	00000580	2**4
	CONTENTS, ALLOC, LOAD, READONLY, CODE					
12	.plt.got	00000008	000000000000005b0	000000000000005b0	000005b0	2**3
	CONTENTS, ALLOC, LOAD, READONLY, CODE					
13	.text	000002d2	000000000000005c0	000000000000005c0	000005c0	2**4
	CONTENTS, ALLOC, LOAD, READONLY, CODE					
14	.fini	00000009	00000000000000894	00000000000000894	00000894	2**2
	CONTENTS, ALLOC, LOAD, READONLY, CODE					
15	.rodata	00000050	000000000000008a0	000000000000008a0	000008a0	2**3
	CONTENTS, ALLOC, LOAD, READONLY, DATA					
...						
31	.debug_str	000002be	00000000000000000	00000000000000000	00001750	2**0
	CONTENTS, READONLY, DEBUGGING					

Some Interesting Sections

Index	Name	Size	Addr	Flags
...				
13	.text	x02d2	5c0	CODE,ALLOC,READONLY
...	...			
15	.rodata	x0050	8a0	DATA,ALLOC,READONLY
...	...			
24	.data	x0010	201028	DATA,ALLOC...
24	.bss	x0008	201038	ALLOC
...				

.text section – x86 binary instructions

```
> objdump -s -j.text x86Math
```

```
x86Math:      file format elf64-x86-64
```

```
Contents of section .text:
```

```
05c0 31ed4989 d15e4889 e24883e4 f050544c 1.I..^H..H...PTL
05d0 8d05ba02 0000488d 0d430200 00488d3d .....H..C...H.=
```

```
...
```

```
> objdump -d -j.text x86Math
```

```
Disassembly of section .text:
```

```
000000000000005c0 <_start>:
5c0:    31  ed                xor     %ebp,%ebp
5c2:    49  89  d1            mov     %rdx,%r9
5c5:    5e                    pop     %rsi
```

```
...
```

.data –initialized global/static data

```
> objdump -s -j.data x86Math
```

```
x86Math:      file format elf64-x86-64
```

```
Contents of section .data:
```

```
201028 00000000 00000000 30102000 00000000  ....0. ....
```

.bss—Size of uninitialized data

- Takes no space in object code (or in file)!
- “bss” acronym for “Block Storage Start”
 - or “Better Save Space”
- Header indicates where this segment should start in memory, and how many bytes should be reserved for uninitialized data

.rodata—Read only data (constants)

```
> objdump -s -j.rodata x86Math
```

```
x86Math:      file format elf64-x86-64
```

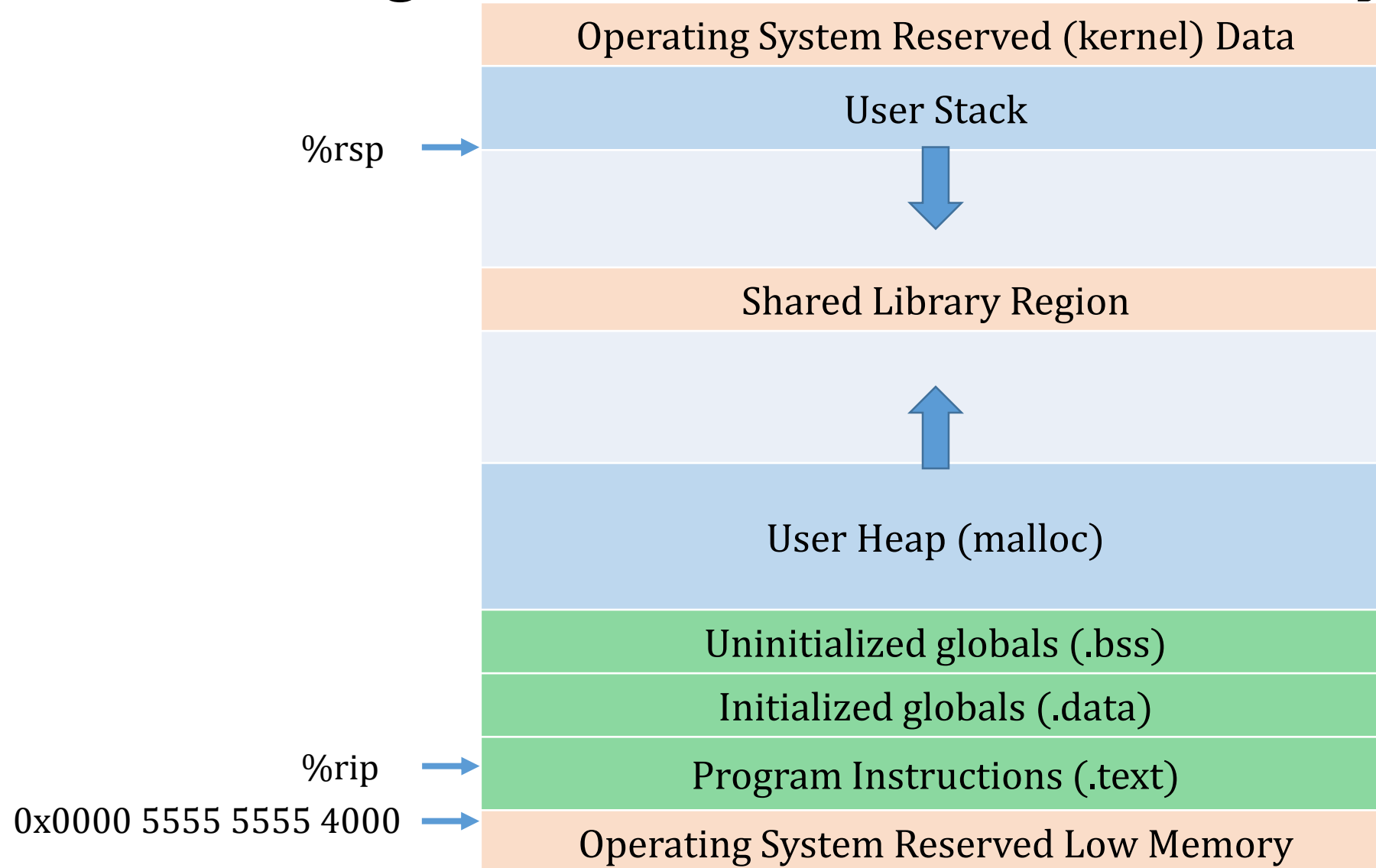
```
Contents of section .rodata:
```

```
08a0 01000200 00000000 496e766f 6b652061 .....Invoke a
08b0 73202573 203c6e75 6d626572 3e0a0000 s %s <number>...
08c0 783d2564 2c207820 73717561 72656420 x=%d, x squared
08d0 2d203478 202b2034 203d2564 20646976 - 4x + 4 =%d div
08e0 69646564 20627920 782d323d 25640a00 ided by x-2=%d..
```

Running a command

- Loading the command:
 - Find the binary file in the file system
 - Create a new address space (and process ID) for the command
 - Load the binary file into memory (new address space)
 - Parse the command line into argc and argv
- Executing the command:
 - Put argc in %rdi and argv in %rsi
 - callq main (address is in symbol table)

Loading an ELF file into Memory



Symbol Table

- ELF keeps track of various “symbol” names ; names of sections, functions, external global variables, etc.
- These are kept in a “symbol table”

Offset	Flags	Section	Length	Name
00000000000000238	l d	.interp	0000000000000000	.interp
000000000000005c0	l d	.text	0000000000000000	.text
000000000000007cf	g F	.text	0000000000000026	add
00000000000000000	F	*UND*	0000000000000000	printf@@GLIBC_2.2.5

- Loader resolves unresolved symbols

Object Code

- It is possible to compile a single file of a multi-file program, and save the x86 code generated
- The result is saved in an ELF format file <pgm>.o
- References to functions outside this file are unresolved, but still in the symbol table
- A “Link Editor” combines object files from several C files into a single executable
 - Invoked by the gcc command
 - Resolves any symbols it can

Make w/ Multiple Compiles

test : fact

./fact

fact: fact.o stack.o

gcc -g -Wall -rdynamic -o fact fact.o stack.o

fact.o : fact.c stack.h

gcc -g -Wall -c fact.c

stack.o : stack.c stack.h

gcc -g -Wall -c stack.c

clean:

-rm fact fact.o stack.o

Link Editor resolves symbols

- `objdump -d fact.o...`

```
82:      e8 00 00 00 00      callq 87 <factorial+0x54>
87:      8b 45 fc           mov    -0x4(%rbp),%eax
```

- `objdump -d fact`

```
b42:      e8 05 00 00 00      callq b4c <printStatsInfo>
b47:      8b 45 fc           mov    -0x4(%rbp),%eax
```

- `gdb --args fact`

```
0x000055555554b42 <+79>:      callq 0x55555554b4c <printStatsInfo>
=> 0x000055555554b47 <+84>:      mov    -0x4(%rbp),%eax
```

Dynamically Resolved Symbols

- Library code often in “Dynamically Loadable Libraries” (DLL’s)
- DLL’s are loaded when they are first used
 - Prevents loading enormous amounts of library code for each program
 - May be shared across different programs
- Invocation of a function in a DLL is resolved when it is called
 - Code invokes a wrapper function, e.g. `printf@plt`
 - `printf@plt` function knows if `printf` has been resolved
 - If not, load the DLL, find the address of `printf`, and save it
 - In either case, call the real `printf`