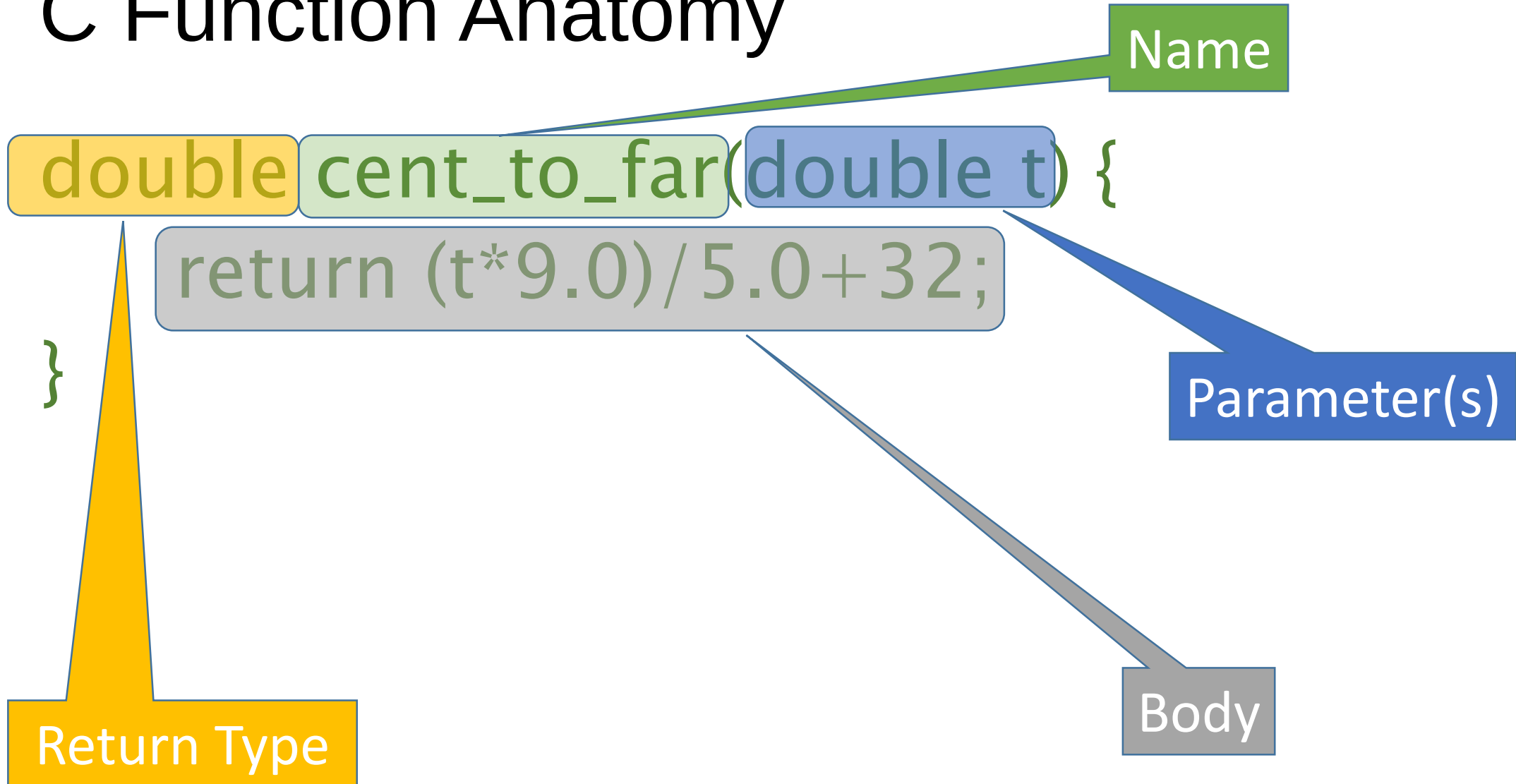


Functions

The Basis of C

C Function Anatomy



Function Invocation

- Invoke a function in a C expression by specifying:

function_name (argument_expression_list)

- *function_name* : name of a previously declared function
- *argument_expression_list* : comma separated list of expressions to be used as arguments
- When a function is invoked, C will....
 1. evaluate argument expressions,
 2. Copy argument values into parameters
 3. invoke the specified function (run it's body)
 4. conceptually replace the function invocation with the returned value

```
float warmer=cent_to_far(ctemp*1.10);
```

Functions for Abstraction

- Function Prototype: `float cent_to_far(float t)`
- Function Behavior: Convert Centigrade to Fahrenheit
- Function Embodiment: WHO CARES?
 - As long as it works, we can ignore the embodiment!
 - Code it once, test it, and then use it lots of times!



The “main” function

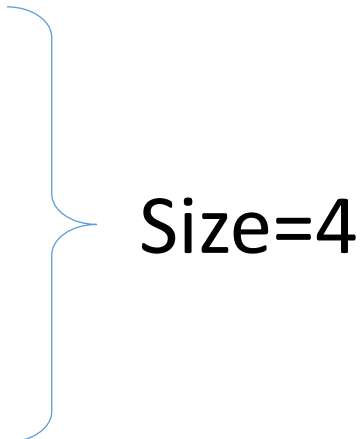
- Every C program must have a “main” function
- When C program is run, the OS invokes the main function
- When the main function returns, program ends
- Return value “int”
 - Return value 0 indicates program worked OK
 - Return value other than 0 indicates program failed
- main function arguments – stay tuned
- main function may invoke lower level functions
- For now: `int main() { ... ; return 0; }`



Arguments to main

- When the operating system calls the “main” function, it:
- parses the command line, splitting at “white space” (blanks, tabs)
 > `./convertTemp 21.3 F C`
- Counts the number of blank delimited words: `argc=4`
- Creates an array of “words” or strings (really an array of arrays)

<code>argv[0]</code>	<code>“./convertTemp”</code>
<code>argv[1]</code>	<code>“21.3”</code>
<code>argv[2]</code>	<code>“F”</code>
<code>argv[3]</code>	<code>“C”</code>



Parameters of main

- Typically use `argc` and `argv` as the names of the parameters to main
`int main(int argc, char **argv) { ...`
- `argc` argument count – the size of the `argv` array
- `argv` argument value (or vector) - a list of strings
 - Each string is an array of characters
 - So `argv` is an array of arrays of characters
 - Almost `(char[])[] argv` – `argv` is an array of (array of characters)
 - (note: not `char[][] argv` – not a two dimensional array of characters!)
 - Can reference each string as `argv[i]`
 - Can reference individual letters as `argv[i][j]` really `(argv[i])[j]`

Function Prototype Declare

- Function Prototype: `float freezingPoint(char scale);`
- Prototype consists of return type, function name, & parameter list
- We can **declare** a function by specifying the prototype;
 - Followed by a semi-colon
- Still need full function **definition** somewhere else
- Enables “Right Side Up” coding
 - Function Prototypes at the top of the file
 - Function definition for top level function (main) next
 - Then function definitions for lower level functions

Function Internals (abstractly)



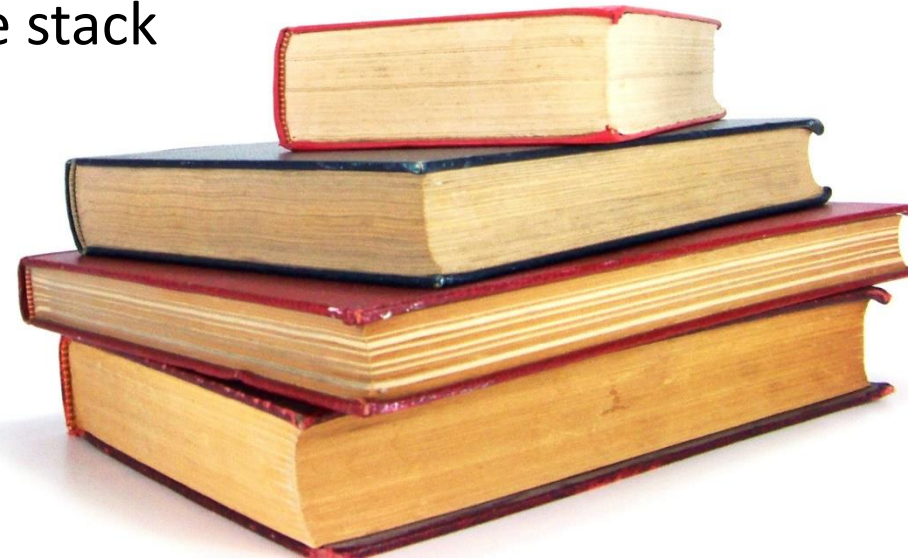
Activation Record

- When a function is called, an “activation record” is created
- Activation records hold:
 - Location in C code where function is invoked and will return
 - Function being invoked
 - Copies of the argument values for this invocation
 - Space for function variable values
 - Space for a return value
- After the function returns, the activation record is deleted

Inv	Fn	args	vars	Ret
main.7	myfn	x=7	a=3,b=4	?

Activation Records are kept in a stack

- “Current” function on top of stack
- When a function is invoked,
 - add its activation record to the top of the stack
- After a function returns,
 - remove it’s activation record from the top of the stack



Example Call Stack

```
1. int addem(int x, int y);  
➔ 2. int main() {  
3.     int a=addem(3,4);  
4.     a=addem(a,4);  
5.     return 0;  
6. }  
7. int addem(int x, int y) { return x+y;}
```

Inv	Fn	args	vars	Ret
OS	main		a=	

Example Call Stack

```
1. int addem(int x, int y);  
2. int main() {  
→ 3.   int a=addem(3,4);  
4.   a=addem(a,4);  
5.   return 0;  
6. }  
7. int addem(int x, int y) { return x+y;}
```

Inv	Fn	args	vars	Ret
3.10	addem	x=3,y=4		

Inv	Fn	args	vars	Ret
OS	main		a=	

Example Call Stack

```
1. int addem(int x, int y);  
2. int main() {  
3.     int a=addem(3,4);  
4.     a=addem(a,4);  
5.     return 0;  
6. }
```

➔ 7. int addem(int x, int y) { return x+y;}

Inv	Fn	args	vars	Ret
3.10	addem	x=3,y=4		7

Inv	Fn	args	vars	Ret
OS	main		a=	

Example Call Stack

```
1. int addem(int x, int y);  
2. int main() {  
→ 3.   int a=addem(3,4);  
4.   a=addem(a,4);  
5.   return 0;  
6. }  
7. int addem(int x, int y) { return x+y;}
```

Inv	Fn	args	vars	Ret
OS	main		a=7	

Example Call Stack

```
1. int addem(int x, int y);  
2. int main() {  
3.     int a=addem(3,4);  
➔ 4.     a=addem(a,4);  
5.     return 0;  
6. }  
7. int addem(int x, int y) { return x+y;}
```

Inv	Fn	args	vars	Ret
4.7	addem	x=7,y=4		

Inv	Fn	args	vars	Ret
OS	main		a=7	

Example Call Stack

```
1. int addem(int x, int y);  
2. int main() {  
3.     int a=addem(3,4);  
4.     a=addem(a,4);  
5.     return 0;  
6. }
```

➔ 7. int addem(int x, int y) { return x+y;}

Inv	Fn	args	vars	Ret
4.7	addem	x=7,y=4		11

Inv	Fn	args	vars	Ret
OS	main		a=7	

Example Call Stack

```
1. int addem(int x, int y);  
2. int main() {  
3.     int a=addem(3,4);  
→ 4.     a=addem(a,4);  
5.     return 0;  
6. }  
7. int addem(int x, int y) { return x+y;}
```

Inv	Fn	args	vars	Ret
OS	main		a=11	

Example Call Stack

```
1. int addem(int x, int y);  
2. int main() {  
3.     int a=addem(3,4);  
4.     a=addem(a,4);  
→ 5.     return 0;  
6. }  
7. int addem(int x, int y) { return x+y;}
```

Inv	Fn	args	vars	Ret
OS	main		a=11	0

Notes on Call Stacks

- Arguments are passed by value
 - Can't update callers value!
- GDB “where” command prints call stack
- Function variable management is useful
 - Don't need to worry about caller's environment!
 - Don't need to worry about previous invocations
 - Enables “recursion”

- A “recursive” function is a function which calls itself
- For example, “factorial”
 - $\text{fact}(1)=1$
 - $\text{fact}(n)=n * \text{fact}(n-1)$ for $n>1$
- For example:
 - $\text{fact}(2) = 2 \times \text{fact}(1) = 2 * 1 = 2$
 - $\text{fact}(3) = 3 \times \text{fact}(2) = 3 * 2 = 6$
 - $\text{fact}(4) = 4 \times \text{fact}(3) = 4 * 6 = 24$
 - ...



Factorials in C

```
1. int fact(int x);  
→ 2. int main() {  
3.     int a = fact(4);  
4.     return 0;  
5. }  
6. int fact(int x) {  
7.     if (x==1) return 1;  
8.     return x*fact(x-1);  
9. }
```

Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```
1. int fact(int x);  
2. int main() {  
→ 3.     int a = fact(4);  
4.     return 0;  
5. }  
6. int fact(int x) {  
7.     if (x==1) return 1;  
8.     return x*fact(x-1);  
9. }
```

Inv	Fn	args	vars	Ret
3.12	fact	x=4		

Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```
1. int fact(int x);  
2. int main() {  
3.     int a = fact(4);  
4.     return 0;  
5. }  
6. int fact(int x) {  
7.     if (x==1) return 1;  
8.     return x*fact(x-1);  
9. }
```



Inv	Fn	args	vars	Ret
8.13	fact	x=3		

Inv	Fn	args	vars	Ret
3.12	fact	x=4		

Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```

1. int fact(int x);
2. int main() {
3.     int a = fact(4);
4.     return 0;
5. }
6. int fact(int x) {
7.     if (x==1) return 1;
8.     return x*fact(x-1);
9. }
    
```



Inv	Fn	args	vars	Ret
8.13	fact	x=2		

Inv	Fn	args	vars	Ret
8.13	fact	x=3		

Inv	Fn	args	vars	Ret
3.12	fact	x=4		

Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```

1. int fact(int x);
2. int main() {
3.     int a = fact(4);
4.     return 0;
5. }
6. int fact(int x) {
7.     if (x==1) return 1;
8.     return x*fact(x-1);
9. }

```



Inv	Fn	args	vars	Ret
8.13	fact	x=1		

Inv	Fn	args	vars	Ret
8.13	fact	x=2		

Inv	Fn	args	vars	Ret
8.13	fact	x=3		

Inv	Fn	args	vars	Ret
3.12	fact	x=4		

Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```

1. int fact(int x);
2. int main() {
3.     int a = fact(4);
4.     return 0;
5. }
6. int fact(int x) {
→ 7.     if (x==1) return 1;
8.     return x*fact(x-1);
9. }
```

Inv	Fn	args	vars	Ret
8.13	fact	x=1		1

Inv	Fn	args	vars	Ret
8.13	fact	x=2		

Inv	Fn	args	vars	Ret
8.13	fact	x=3		

Inv	Fn	args	vars	Ret
3.12	fact	x=4		

Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```

1. int fact(int x);
2. int main() {
3.     int a = fact(4);
4.     return 0;
5. }
6. int fact(int x) {
7.     if (x==1) return 1;
8.     return x*fact(x-1);
9. }

```

1

Inv	Fn	args	vars	Ret
8.13	fact	x=2		2

Inv	Fn	args	vars	Ret
8.13	fact	x=3		

Inv	Fn	args	vars	Ret
3.12	fact	x=4		

Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```

1. int fact(int x);
2. int main() {
3.     int a = fact(4);
4.     return 0;
5. }
6. int fact(int x) {
7.     if (x==1) return 1;
8.     return x*fact(x-1);
9. }
    
```

2

Inv	Fn	args	vars	Ret
8.13	fact	x=3		6

Inv	Fn	args	vars	Ret
3.12	fact	x=4		

Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```

1. int fact(int x);
2. int main() {
3.     int a = fact(4);
4.     return 0;
5. }
6. int fact(int x) {
7.     if (x==1) return 1;
8.     return x*fact(x-1);
9. }
    
```

6

Inv	Fn	args	vars	Ret
3.8	fact	x=4		24
Inv	Fn	args	vars	Ret
OS	main		a	

Factorials in C

```

1. int fact(int x);
2. int main() {
3.     int a = fact(4);
4.     return 0;
5. }
6. int fact(int x) {
7.     if (x==1) return 1;
8.     return x*fact(x-1);
9. }
    
```

24

Inv	Fn	args	vars	Ret
OS	main		a=24	

Notes on Recursion

- Recursive functions always recurse to a **simpler** case
 - In “fact” – we always decrease argument by 1
- Recursive functions always have a “base” – a **simplest** case
 - Must be non-recursive to stop the recursion!
 - A condition that stops the recursion
 - In “fact” – base is the “1” case
- Recursive functions are very similar to inductive proofs in math
 - If the recursion works for 1, and if you assume the recursion works for $n-1$ and show it works for n , then the recursion works for any integer
 - $\text{fact}(1)=1$, and If $\text{fact}(3)=3 \times 2 \times 1$, then $\text{fact}(4) = 4 \times \text{fact}(3) = 4 \times 3 \times 2 \times 1$

Recursion can be replaced by loops

```
1. int fact(int x);  
2. int main() {  
3.     int a = fact(4);  
4.     return 0;  
5. }  
6. int fact(int x) {  
7.     int f=1; int j; for(j=2;j<x;j++) f*=j;  
8.     return f;  
9. }
```

Recursion vs. Loops

- Recursion is often conceptually very simple, loops can be more complicated to understand
- Loops often seem to be significantly more efficient
 - Don't have the overhead of function invocation, creating call stacks, etc.
- However, recursion is so common, the compiler optimizes recursion, often turning it into a loop

Bottom line: Use whatever makes the most sense to you, and let the compiler take care of efficiency

Resources

- Programming in C, Chapter 7
- YouTube: Meat-a-Morphis – Introduction to Functions
<https://www.youtube.com/watch?v=VUTXsPFx-qQ>
- Wikipedia: Subroutine <https://en.wikipedia.org/wiki/Subroutine>
- Wikipedia: C Standard Library:
https://en.wikipedia.org/wiki/C_standard_library
- Wikipedia: Recursion (computer science)
[https://en.wikipedia.org/wiki/Recursion_\(computer_science\)](https://en.wikipedia.org/wiki/Recursion_(computer_science))
- C-FAQ Library Functions: <http://c-faq.com/lib/index.html>
- Linux MAN pages online (library functions): http://man7.org/linux/man-pages/dir_all_alphabetic.html