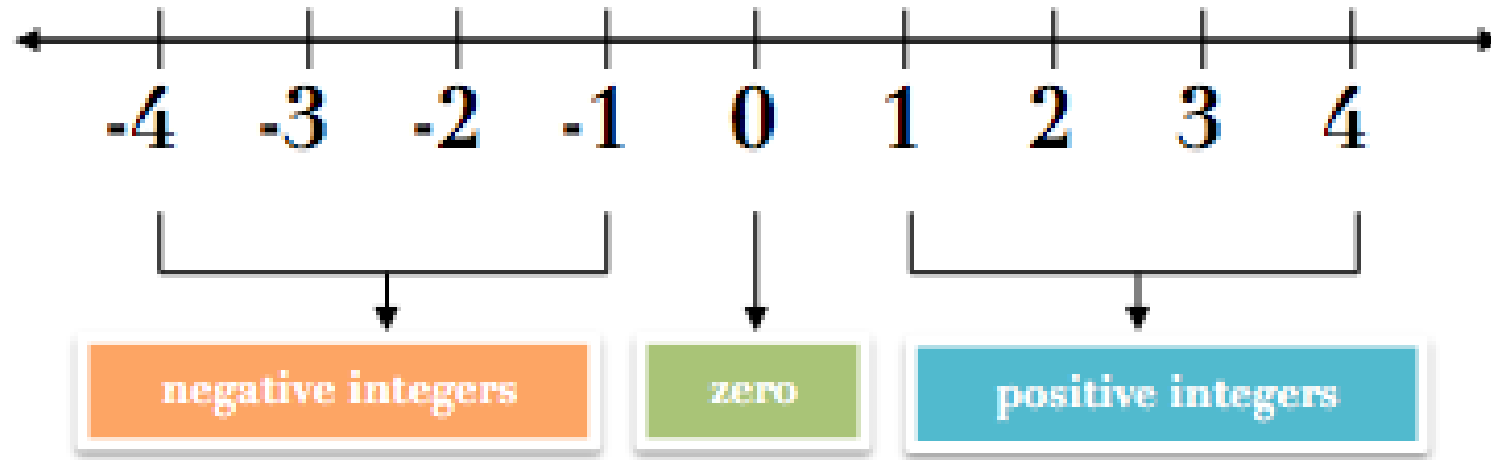
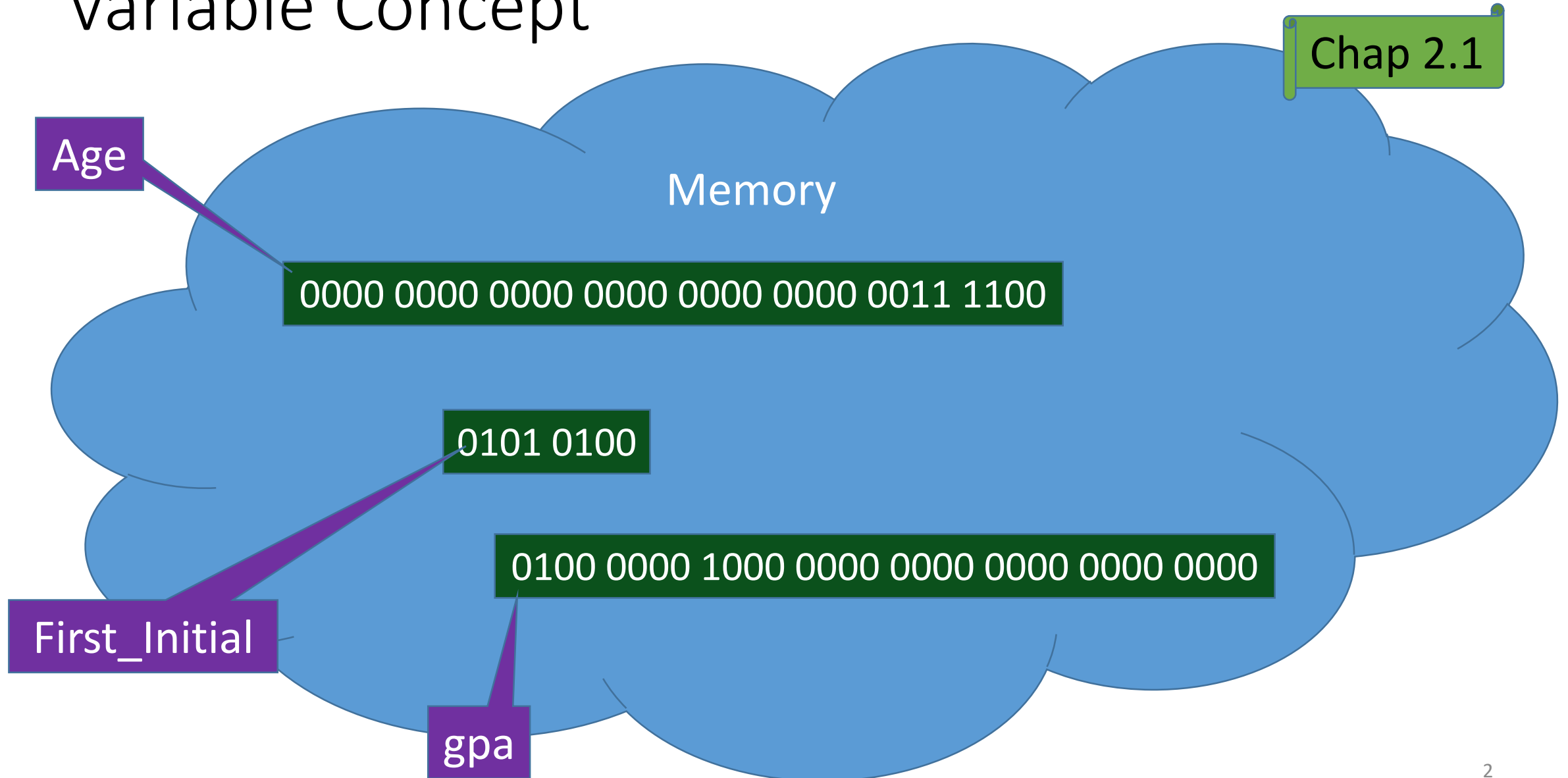




Data in C: Integers

Variable Concept

Chap 2.1



Data Types

- Tag each piece of data / location in memory with a “type”
- Type tells compiler how many bits to use
- Type tells compiler how to interpret those bits
- Enables automatic conversion from one type to another
- Enables compiler to check to make sure data is used correctly

C Built-in Types

Numbers

Text

void

Integer
Binary 2's complement

Real
IEEE 754

char
ASCII

char
8 bit

short
16 bit

int
32 bit

long
64 bit

float
32 bit

double
64 bit

Two's Complement Binary

Chap 2.2

- Number represented by a vector of n binary digits (1's or 0's)
- If leftmost bit is a zero, number is positive, simple binary

$$value = \sum_{i=0}^{n-1} b_i \times 2^i$$

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
0	0	0	0	1	0	1	0

- If leftmost bit is a one, number is negative

$$value = \left(\sum_{i=0}^{n-1} b_i \times 2^i \right) - 2^n$$

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
1	1	1	1	1	0	1	0

C Idiom: (Using Leaks) 2^x

- In binary representation, 2^x is just a single 1 in the x th bit
 - e.g. $2^3 = 0b0000\ 1000 = 8$
- Use the “shift” operator, $<<$, to calculate 2^x

```
int x=3;  
int y=1<<x;  
printf("y is %d\n",y); // prints y is 8
```

C idiom: (using Leaks) Strings of x 1 bits

- Use the fact that $2^x - 1$ is a string of x 1 bits
 - e.g. $2^3 = 8 = 0b000\ 1000$
 $-1 = 0b000\ 0001$

• $2^3 - 1 = 7 = 0b000\ 0111$
- Combine 2^x idiom with -1

```
int x=3;  
int y=(1<<x)-1; // y is 0b0000 ... 0000 0111
```

Two's Complement Leaky Shortcut

- To multiply by -1, flip the bits and add 1
 - “Flip this bits” is the same as $0b111\dots 1111 - u_value$
 - $1-0=1$, flip a zero to a one
 - $1-1=0$, flip a one to a zero
 - Never carry
 - $0b111\dots 1111$ is 2^n-1 , so “flip the bits” is $(2^n-1)-u_value$
 - $u_value = \sum_{i=0}^{n-1} b_i \times 2^i$
 - flip the bits and add 1 is:

$$\begin{aligned} (2^n-1)-u_value+1 &= 2^n-u_value \\ &= -(u_value-2^n) \\ &= -value \end{aligned}$$

C idiom: Bitwise AND to select bits

- Create a “mask” that has 1 bits for every interesting bit, and 0 for every bit you want to ignore
- Bitwise AND the mask with the original value
- Result has the original bit in the interesting positions
 - $0 \& 1 = 0$, $1 \& 1 = 1$
- Result has 0 in the ignorable positions
 - $0 \& 0 = 0$, $1 \& 0 = 0$
- Evaluate “truth” of result – “True” means at least one interesting bit was on (1), “False” means all interesting bits were off (0).

Example: Print Binary...

```
#include <stdio.h>
#include <stdlib.h>
int main(int argc, char **argv) {
    int n=atoi(argv[1]); int i; printf("%d = 0b",n);
    for(i=31;i>=0;i--) {
        printf("%c",(n&1 <<i)?'1':'0');
        if (0==i%4) printf(" ");
    }
    printf("\n"); return 0;
}
```

Using Hexadecimal Numbers

- Shorthand for writing bits
- Each hexadecimal digit represents 4 bits
- Much easier to read/write/remember

Dec	Hex	Bin
0	0x0	0b0000
1	0x1	0b0001
2	0x2	0b0010
3	0x3	0b0011
4	0x4	0b0100
5	0x5	0b0101
6	0x6	0b0110
7	0x7	0b0111
8	0x8	0b1000
9	0x9	0b1001
10	0xA	0b1010
11	0xB	0b1011
12	0xC	0b1100
13	0xD	0b1101
14	0xE	0b1110
15	0xF	0b1111

Example: Print hex (without %x)

```
int main(int argc, char **argv) {
    long n=atol(argv[1]);
    int i; printf("%ld = 0x",n);
    unsigned long mask=(1<<4)-1; // Four ones in the rightmost byte
    char * xd="0123456789ABCDEF";
    for(i=2*sizeof(n)-1;i>=0;i--) {
        int v=(n&(mask<<(i*4)))>>(i*4);
        printf("%c",xd[v]);
        if (0==i%4) printf(" ");
    }
    printf("\n"); return 0;
}
```

Example: alternate print hex (without %x)

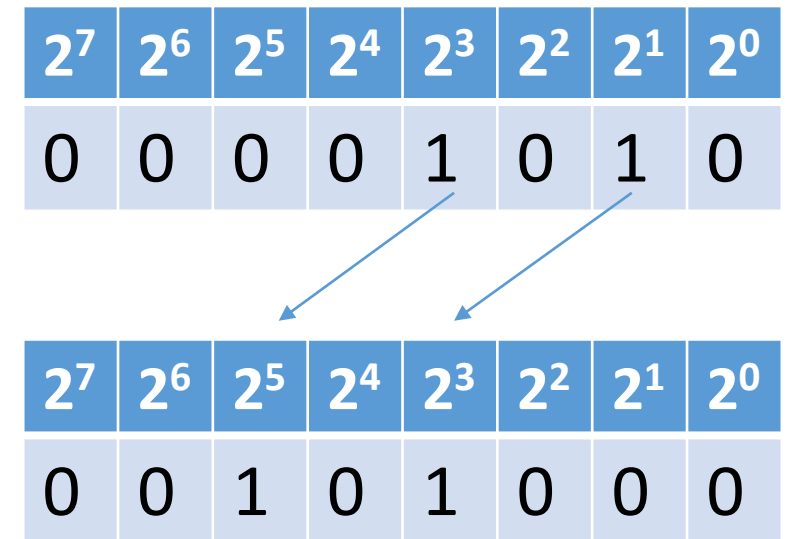
```
int main(int argc, char **argv) {  
    union { long ln; unsigned char lc[8]; } n;  
    char * xd="0123456789ABCDEF";  
    n.ln=atol(argv[1]); printf("%ld = 0x",n.ln);  
    for(int i=0;i<sizeof(n);i++) {  
        if (i>0 && 0==i%2) printf(" ");  
        printf("%c%c",xd[(n.lc[i]&0xF0)>>4],xd[n.lc[i]&0x0F]);  
    }  
    printf("\n");  
    return 0;  
}
```

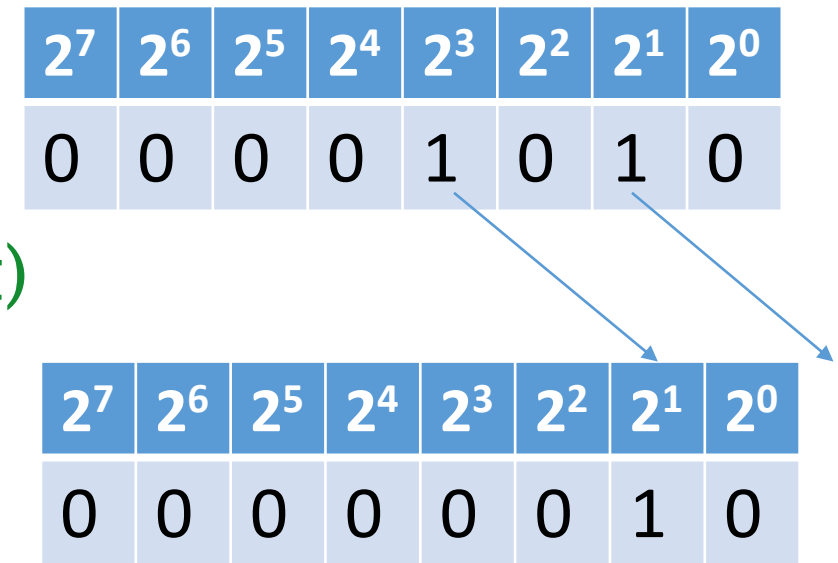
C Leaky Idiom – Multiply by 2^x

- Instead of multiplying a number by a power of two, shift it to the left

```
int x=2;  
int y=atoi(argv[1]);  
int ym = y<<x; // * ym = y * 2^x
```

- Shifts are faster than multiplication





Unsigned Integers

Chap 2.2.2

- C has an “unsigned” keyword that can be used to modify type
 - e.g. `unsigned short x;`
- Enables representation in raw binary – not two’s complement
- Extra bit enables two-times larger maximum number
 - Not much of a benefit
 - Impossible to represent negative numbers
- Unsigned often leads to counter-intuitive results!
 - More explanation coming later



Q: Why four flavors of Integers?

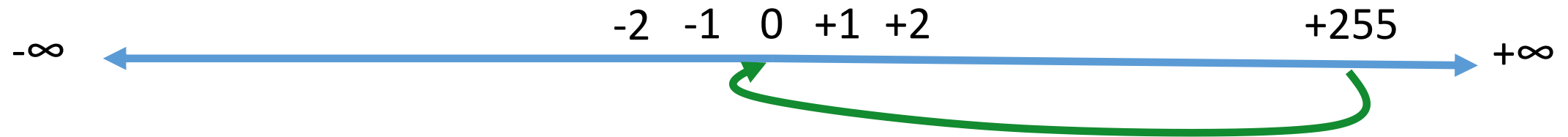
- A: Space vs. Capacity trade-off

Type	# Bits ¹	Min	Max	U Max
char	8	$-2^7 = -128$	$2^7 - 1 = 127$	$2^8 - 1 = 255$
short	16	$-2^{15} = -32,768$	$2^{15} - 1 = 32,767$	$2^{16} - 1 = 65,535$
int	32	$\sim -2.15 \times 10^9$	$\sim 2.15 \times 10^9$	$\sim 4.3 \times 10^9$
long	64	$\sim -9 \times 10^{18}$	$\sim 9 \times 10^{18}$	$\sim 18.5 \times 10^{18}$

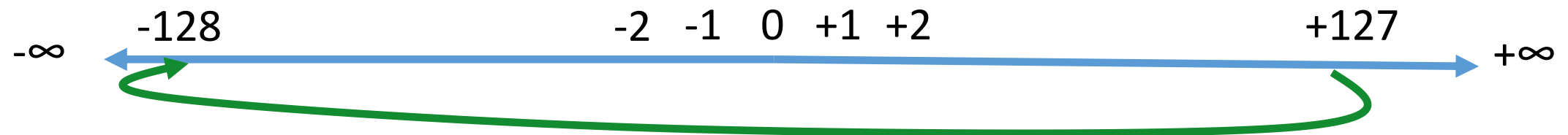
¹Number of bits may be different on different machines!

Leaky Abstraction: Infinite Integers

- In mathematics, every integer, n has a successor, $n+1$
- In C, integers are finite...
 - No indication of overflow!
- For unsigned char, $255 = 0b1111\ 1111 + 1 = 0b0000\ 0000 = 0$



- For char, $127 = 0b0111\ 1111 + 1 = 0b1000\ 0000 = -128$



Constants

- Decimal numbers are signed integers: `int x=93;`
- Add “U” suffix to make them unsigned: `unsigned char uy=183U;`
 - Almost never used – only needed if signed value doesn’t fit
- Prefix with 0 to interpret as an octal number: `int xa=013; // == 11`
- Prefix with x, 0x, X, or 0X for hexadecimal: `int xb=0x12; // == 18`
- Add “L” suffix to make them long : `long yl = 0X000011110000FFFFL;`
 - Almost never used – only needed if constant value doesn’t fit int

Casting : C Data Format Conversion

- When you cast a movie, you decide what role each actor will play
- In programming, we use the term “cast” to indicate what role (type) we expect the bits to play
- In other words, how should your program interpret the specific bits in memory
- Unless explicitly cast, determine the interpretation based on the declared type of a variable



Mixing Integer Types

- Mixed Type Expressions

```
int x; long y; y=y*x;
```

- Assignment Statements

```
int x; long y; x=y*3.0;
```

- Argument Evaluation

```
int myfn(long x); int y=myfn(3);
```

- Explicit Casting

```
int x=7; long y = ((long)x)/3;
```

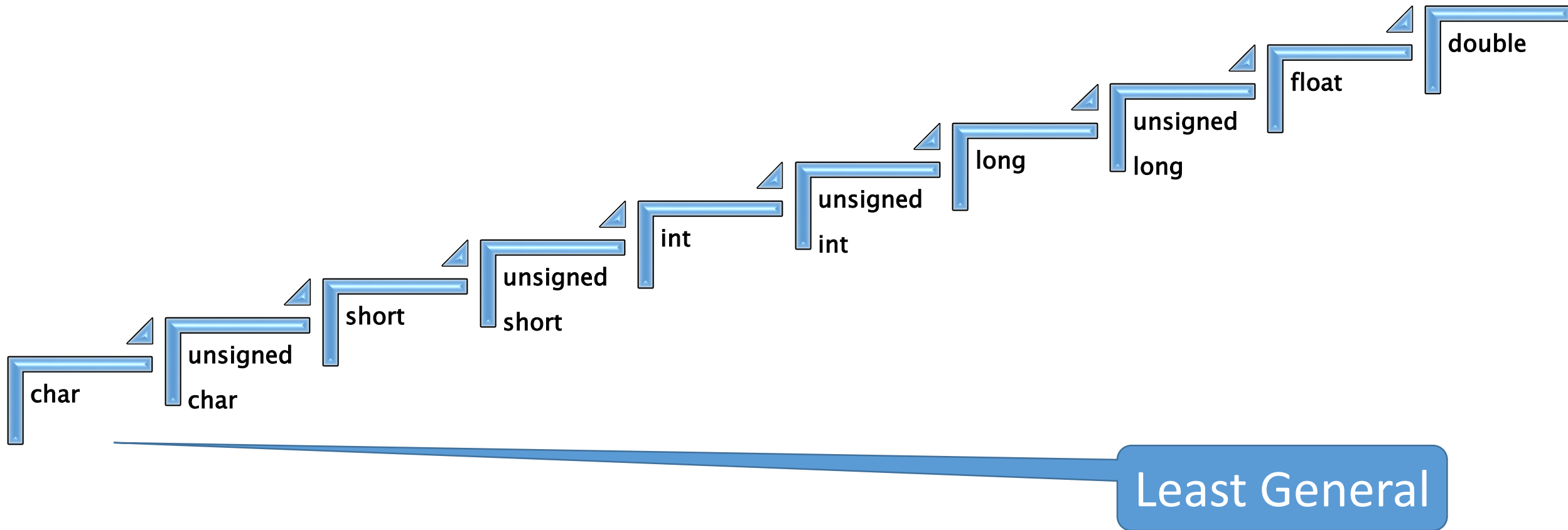
- C compiler could issue an error message for mixed types
- C compiler could require explicit casting for all conversions
- C could automatically convert values, based on pre-defined rules

C Automatic type conversion rules

- In an expression (or part of an expression), C converts all components in that expression to the most “general” type, and then evaluates the expression using that general type
- In an assignment (or argument evaluation), C converts the value of the expression to the type of the receiver
- C converts expressions with a valid explicit cast

Generality of Numeric Types

Most General



Converting Signed vs Unsigned

Chap 2.2.5

- Bits stay exactly the same, but the bits are INTERPRETTED differently

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
1	1	1	0	1	0	1	0

```
char x = -22;  
unsigned char y=x;  
printf("y=%d\n",y);  
// prints y=234
```

```
unsigned char w = 234;  
char z=w;  
printf("z=%d\n",z);  
// prints z=-22
```

Changing Integer Size

Chap 2.2.6

- Truncate or Pad on left with sign bit

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
1	1	1	0	1	0	1	0

2^{15}	2^{14}	2^{13}	2^{12}	2^{11}	2^{10}	2^9	2^8	2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
1	1	1	1	1	1	1	1	1	1	1	0	1	0	1	0

```
char x = -22;
short y=x;
```

```
short w = -22;
char z=w;
```

Changing Integer Size (unsigned)

- Truncate or Pad on left with sign bit (=0)

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
1	1	1	0	1	0	1	0

2^{15}	2^{14}	2^{13}	2^{12}	2^{11}	2^{10}	2^9	2^8	2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
0	0	0	0	0	0	0	0	1	1	1	0	1	0	1	0

```
unsigned char x = 234;  
unsigned short y=x;
```

```
unsigned short w = 234;  
unsigned char z=w;
```

Leaky Abstractions: Conversion Errors

- When C converts a negative signed number to a positive number

```
char x = -1; unsigned char y = x; printf("y = %d\n",y);
```

- When C converts a wide number to a smaller width, but the number doesn't fit

```
short x=260; char y = x; printf("y = %d\n",y);
```

Unsigned Pitfall

```
unsigned int width = +8;  
signed int leftX = -13;  
if ((leftX < 0) && (leftX + width) > 0) {  
    printf("Rectangle crosses y axis\n");  
}
```

Abstraction

Bits are stored in memory from most significant at left
to least significant at right
(int 100,000 = 0x0001 86A0)

2^{31}	2^{30}	2^{29}	2^{28}	2^{27}	2^{26}	2^{25}	2^{24}	2^{23}	2^{22}	2^{21}	2^{20}	2^{19}	2^{18}	2^{17}	2^{16}	2^{15}	2^{14}	2^{13}	2^{12}	2^{11}	2^{10}	2^9	2^8	2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0	1	1	0	1	0	1	0	0	0	0	0
0				0				0				1				8				6				A				0			
Byte m								Byte m+1								Byte m+2								Byte m+3							

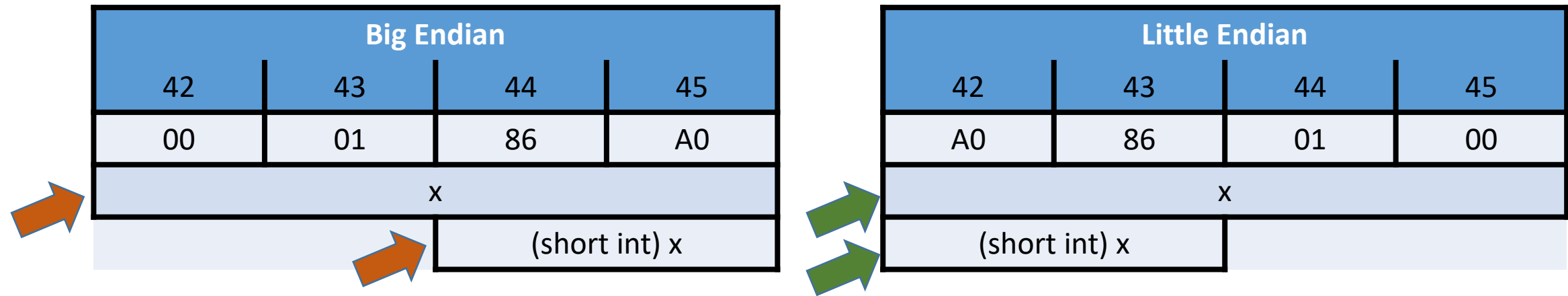
Chap 2.1.3

Chap 2.1.3

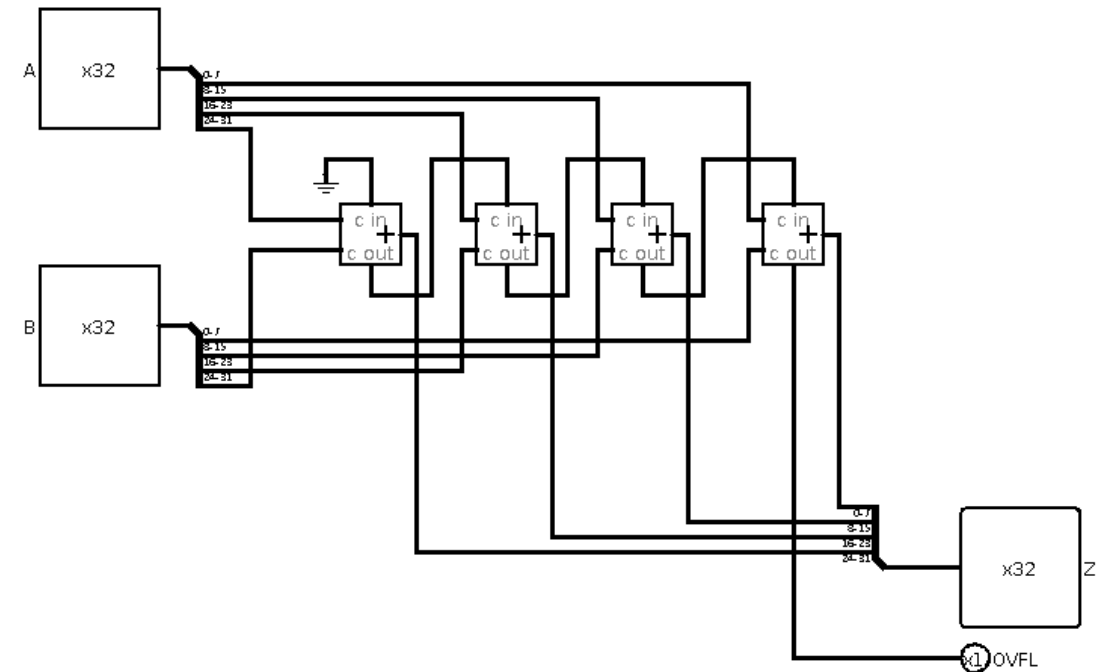
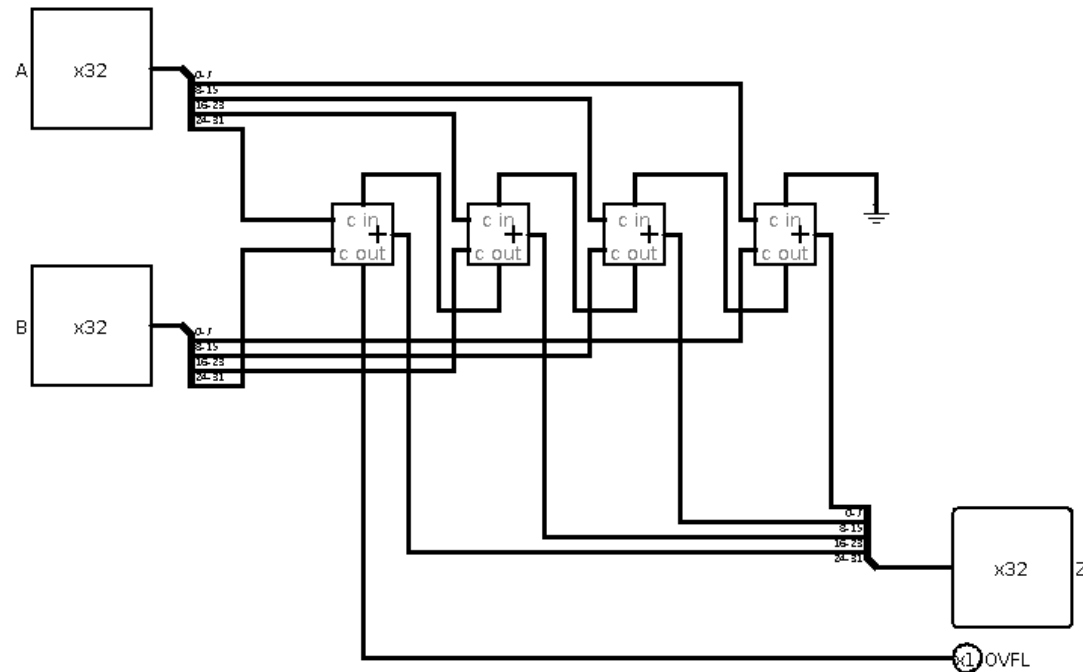
2 ⁷	2 ⁶	2 ⁵	2 ⁴	2 ³	2 ²	2 ¹	2 ⁰	2 ¹⁵	2 ¹⁴	2 ¹³	2 ¹²	2 ¹¹	2 ¹⁰	2 ⁹	2 ⁸	2 ²³	2 ²²	2 ²¹	2 ²⁰	2 ¹⁹	2 ¹⁸	2 ¹⁷	2 ¹⁶	S	2 ³⁰	2 ²⁹	2 ²⁸	2 ²⁷	2 ²⁶	2 ²⁵	2 ²⁴
b ₃₁	b ₃₀	b ₂₉	b ₂₈	b ₂₇	b ₂₆	b ₂₅	b ₂₄	b ₂₃	b ₂₂	b ₂₁	b ₂₀	b ₁₉	b ₁₈	b ₁₇	b ₁₆	b ₁₅	b ₁₄	b ₁₃	b ₁₂	b ₁₁	b ₁₀	b ₉	b ₈	b ₇	b ₆	b ₅	b ₄	b ₃	b ₂	b ₁	b ₀
1	0	1	0	0	0	0	0	1	0	0	0	0	1	1	0	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	0
A				0				8				6				0				1				0				0			
Byte 42								Byte 43								Byte 44								Byte 45							

Why Little Endian?

- Casting/Conversion:
 - `int x; /* 32 bits starting at byte 42 */`
 - `y = (short int) x; /* Put the least significant 16 bits from x into y */`



Big Endian vs. Little Endian Adder



When does Endian-ness Leak?

- Only on multi-byte data types (short, int, long, float, double)
- Only when we look at the underlying bits or transfer binary data
- Big-endian machine: First byte is the most significant byte
 - Everything works until we get binary data from a little-endian machine
- Little-endian machine: First byte is the least significant byte
 - When hardware interprets the bits/bytes as a number, bytes are switched
 - We don't even know if a machine is big-endian or little-endian!
 - Until: we get binary data from a big-endian machine OR
 - Until we look at the bit representation of the data, not treated as a number

Managing Endian-Ness

- Network standard is big-endian
- stdlib functions
 - machine representation → network (big-endian) representation
 - htons (short) , htonl (long)
 - Network representation (big-endian) → machine representation
 - ntohs (short), ntohl (long)
 - No-ops when hardware is big-endian
- endian.h functions
 - htobe16, htobe32, htobe64, htole16, htole32, htole64
 - be16toh, be32toh, be64toh, le16toh, le32toh, le64toh

See xmp_endian/network.c

Endian-ness in this Course

- LDAP machines are little-endian
- When you put **binary** data in a file, you must enter numbers in little-endian format!
- When you print memory from gdb,
 - if you specify byte output formats, data will appear little endian.
 - (When you specify multi-byte numeric output formats, the infrastructure switches the bytes for you.)
- If you make a union of byte-data and multi-byte formats, the byte data will show endian-ness
- Otherwise, everything looks like big-endian!

Resources

- The C Programming Language, (K&R) Sections 2.2 and 2.7
- Wikipedia Type Conversion:
https://en.wikipedia.org/wiki/Type_conversion
- C Tutorial – Cast operator:
<http://www.crasseux.com/books/ctutorial/The-cast-operator.html#The%20cast%20operator>
- C-FaQ Cast Operator Section:
<http://c-faq.com/~scs/cclass/int/sx4bb.html>