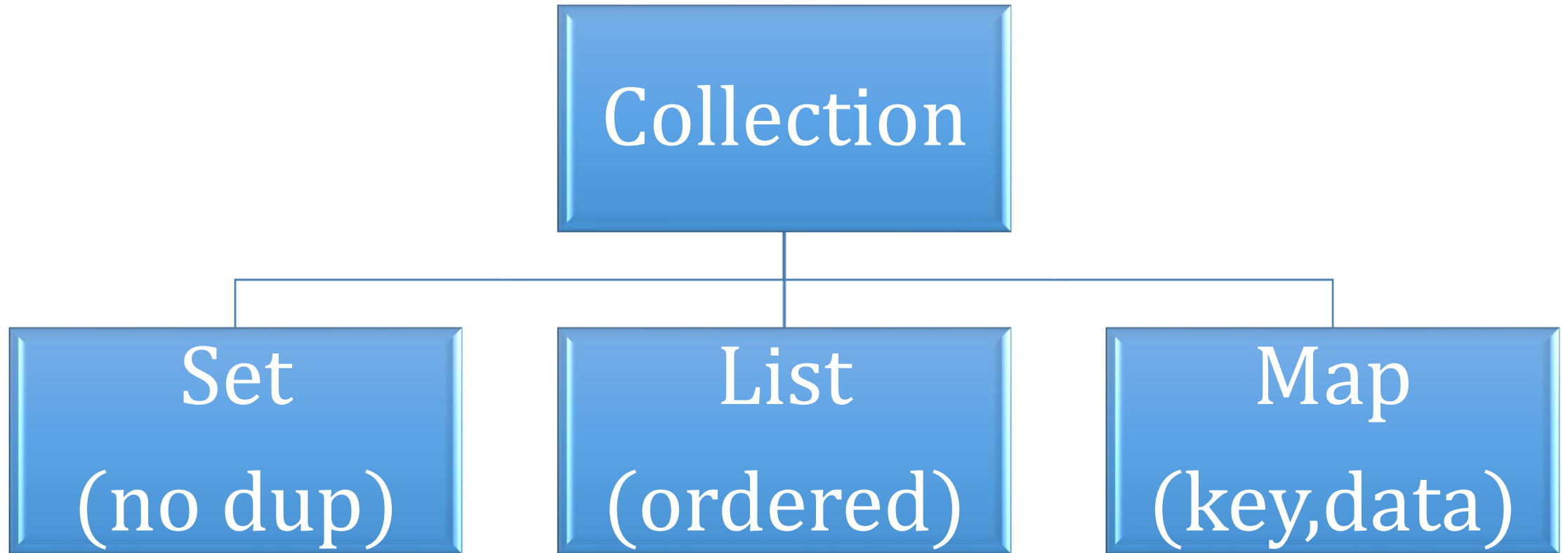


# Using Collections

# The Java Collections Infrastructure





# Typical Collection Usage

- All you need is (for instance) the List interface
- Find a concrete class which implements List (e.g. ArrayList)
- Declare field/variable as List
- Instantiate field/variable using concrete class

```
private List blockList;
```

```
...
```

```
blockList = new ArrayList<Block>();
```

To use a different implementation,  
change this line

# Joel Spolsky\*: Law of Leaky Abstractions

- We like to think of the world abstractly
  - My program uses a list – an ordered collection of elements
- Sometimes we need to know about the concrete implementation
  - If the backing store of a list is an array, insertion/deletion can be slow
  - If the backing store of a list is a linked list, direct indexing can be slow
  - Both act as lists, but one does some list things better than the other
  - We may want to choose an implementation based on our application

\*Author of *Joel on Software* blog, co-founded Stack Overflow

# Some Classes implementing Set

- EnumSet – Backing store: bit vector
  - Requires small fixed enumerated domain
  - Very fast add, remove, contains (one cycle)
  - + methods: allOf(t) clone() complementOf(s) copyOf(c) noneOf(t) of(...e) range(from,to)
- HashSet - Backing store: HashMap
  - constant time add, remove, contains, and size
  - Slow traversal
- LinkedHashSet – Backing Store: HashMap + linked list
  - stabilizes "order" of the set
- TreeSet – Backing Store: TreeMap
  - $\log(n)$  time add, remove, and contains

# Some Classes implementing List

- ArrayList – Backing store: array
  - Fast direct access to elements
  - Occasional slow add/delete to enlarge/shrink array
  - Slower insert/delete from beginning of list
- LinkedList - Backing store: Doubly linked list of nodes
  - constant time add, remove
  - Slow direct access – iterate is faster
  - Also implements Deque and Queue interfaces
- Vector – Backing Store: array-like, but with shrink and grow
  - Thread safe, but slower than ArrayList

# ArrayList vs. array

Function	array	ArrayList
Declare	<i>type[] var</i>	<i>ArrayList&lt;type&gt; var</i> (type must extend Object)
Instantiate	<i>new type[size]</i>	<i>new ArrayList&lt;type&gt;()</i>
Read element i	<i>var[i]</i>	<i>var.get(i)</i>
Write element i	<i>var[i] = value</i>	<i>var.set(i, value)</i>
add element at end	---	<i>var.add(value)</i>
all element in the middle	---	<i>var.add(i, value)</i>
Shrink/Enlarge	instantiate new larger/smaller array and copy old to new	Automatic
Enhanced loop	<i>for(type v: var) { }</i>	<i>for(type v: var) { }</i>
Performance	Good	Equal except inserting early and when grow or shrink is needed

# Some Classes Implementing Map

- EnumMap – Backing Store: array
  - Requires small fixed enumerated key domain
  - Very fast
- HashMap – backing store – array?
  - Very fast insert/delete
  - Slow/unstable traversal
- Hashtable – Thread safe hash map
- LinkedHashMap – backing store: hash map with linked list
  - Stabilizes and speeds up traversal
- TreeMap
- WeakHashMap



# Binning for Unsorted Items

- Keep two or more bins... lists of objects... bins are a list of lists
- Quick function to determine what bin an element belongs in
- Trick is to equalize binsize... so for  $m$  bins, binsize  $\sim n/m$
- Time to insert : find bin, add to bin - fast
- Time to search : find bin, search in bin –  $O(n/m)$
- Time to delete : find bin, find in bin, delete –  $O(n/m)$
- More bins mean faster access, but more memory



# Ultimate binning

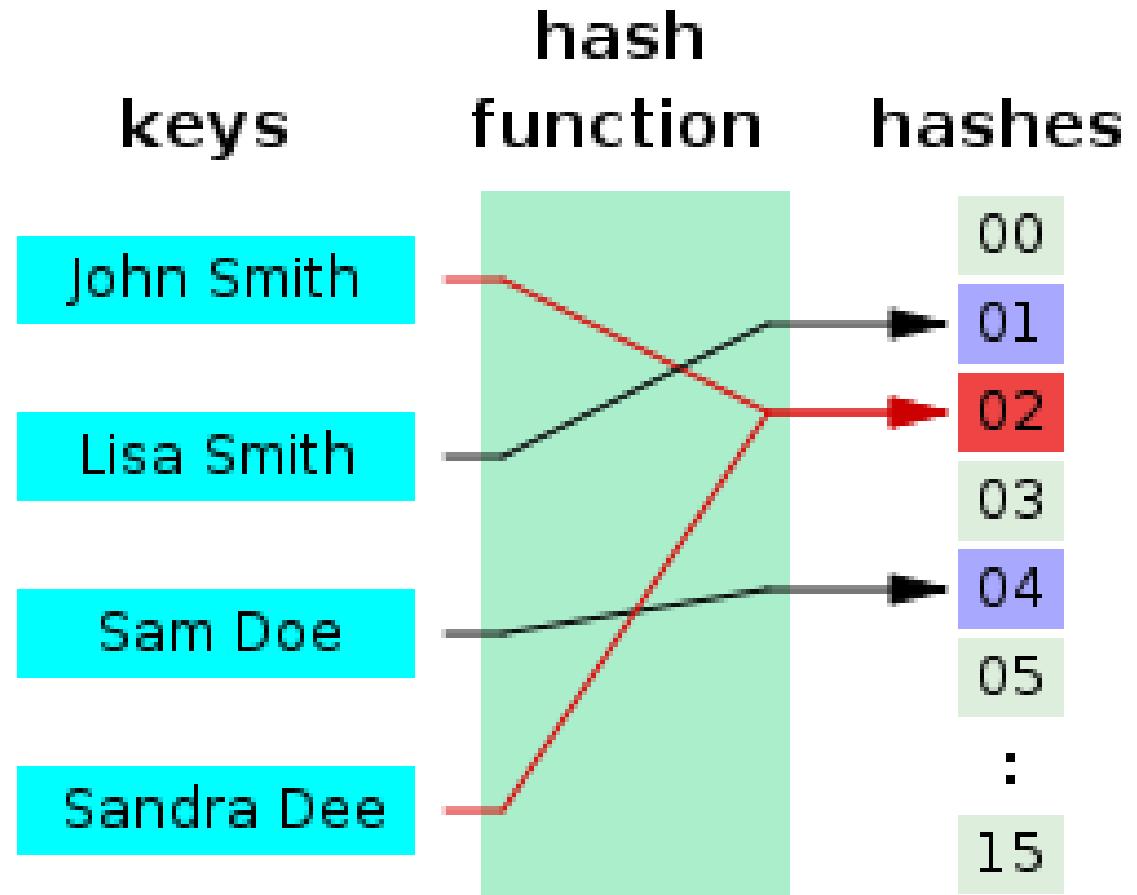
- Separate bin for each item
- Problem... need a specialized function to determine what bin element  $x$  is in
  - Needs to run fast
  - Needs to guarantee that if two elements are the same, they go to the same bin
  - Needs to guarantee that two different elements go to different bins
- Problem: Sparse usage... most bins remain unused
  - Consider a bin for each Lottery number – e.g. pick 6: 45 55 32 91 40 46
  - There are  $10^{12}$  possible lottery tickets!

# Hashing

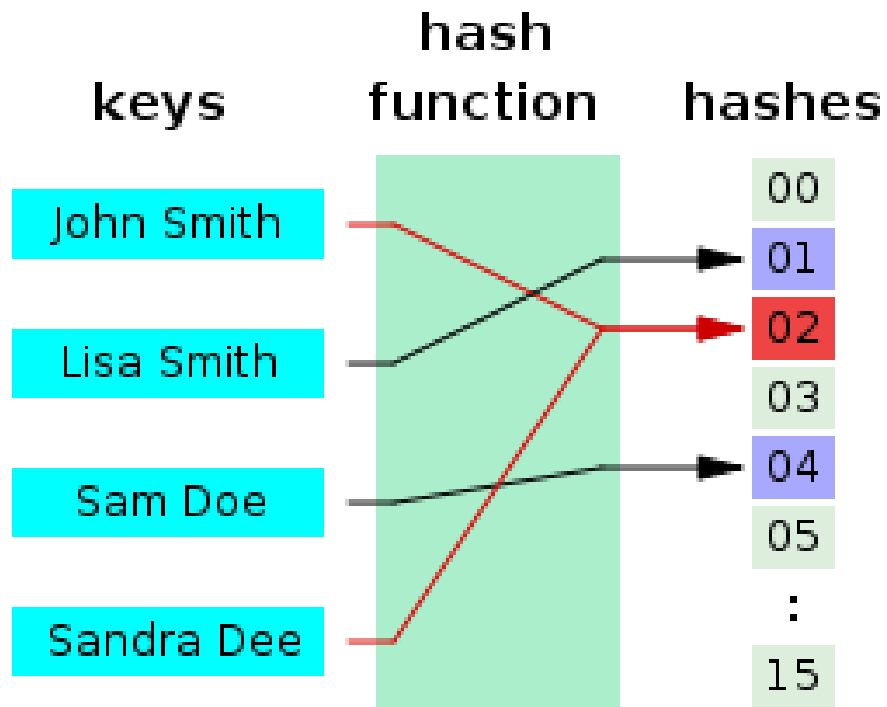
- Hashing is a form of a binning algorithm
- Function to translate from “key” to an index in an array
- Number of bins = number of elements in the array
- The function that performs the translation is a “hash” function
$$index = hash(key)$$
- Guarantee: if  $key_1 == key_2$ , then  $hash(key_1) == hash(key_2)$
- Not guaranteed: if  $key_1 \neq key_2$ , then  $hash(key_1) = ? hash(key_2)$
- “Hash Collision” if  $key_1 \neq key_2$ , but  $hash(key_1) == hash(key_2)$
- Hash function designed to minimize collisions

# Example Hash

- Translate keys to index 0-15
- Each key hashes to the same index every time
- Multiple keys may map to a single index



# Example Hash Table



	Name	Town	ID
0			
1	Lisa Smith	Vestal	6894
2	John Smith	Endicott	1548
	Sandra Dee	Binghamton	6442
3			
4	Sam Doe	Johnson City	2954
5			
...			
15			



# Warning: Modifying Collections in loops

```
for ( Block b : blockList) {  
    if (!b.isUsed()) blocklist.remove(b);  
}
```

- Question: do any of these work?

```
/* ALTERNATE... */  
for (int i=0; i<blocklist.size();i++) {  
    if (!blocklist.get(i).isUsed())  
        blocklist.remove(i);  
}
```

```
/* ALTERNATE 2 ... */  
for ( Iterator it=blocklist.iterator(); it.hasNext();) {  
    Block b = it.next();  
    if (!b.isUsed()) it.remove();  
}
```