

Parallel Processing of Large-Scale XML-Based Application Documents on Multi-core Architectures with PiXiMaL

Michael R. Head
Madhusudhan Govindaraju

Department of Computer Science
Grid Computing Research Laboratory
Binghamton University
`mike@cs.binghamton.edu`
`mgovinda@cs.binghamton.edu`

December 7-12, 2008

Outline

1 Introduction and Motivation

- XML and SOAP
- Ubiquity of Multi-processing Capabilities

2 Related Work

- High Performance XML Processing Approaches

3 Work Completed

- PIXIMAL: Parallel Approach for Processing XML

XML Defined

- Text based (usually UTF-8 encoded)
- Tree structured
- Language independent
- Generalized data format

Motivation from SOAP

- Generalized RPC mechanism (supports other models, too)
- Broad industrial support
- Web Services on the Grid
 - OGSA: Open Grid Services Architecture
 - WSRF: Web Services Resource Framework
- At bottom, SOAP depends on XML

XML Exclusive of SOAP

- General structured data format
- Becoming standard for many scientific datasets
 - HapMap - mapping genes
 - Protein Sequencing
 - NASA astronomical data
 - Many more instances

Explosion of Data

- Enormous increase in data from sensors, satellites, experiments, and simulations*
- Use of XML to store these data is also on the rise
- XML is in use in ways it was never really intended (GB and large size files)

Prevalence of Parallel Machines

- All new high end and mid range CPUs for desktop- and laptop-class computers have at least two cores
- The future of AMD and Intel performance lies in increases in the number of cores
- Despite extant SMP machines, many classes of software applications remain single threaded
 - Multi-threaded programming considered “hard”
 - Reinforced in the current curricula and by existing languages and tools

XML and Multi-Core

- Most string parsing techniques rely on a serial scanning process
- **Challenge:** Existing (singly-threaded) XML parsers are already very efficient [Zhang et al 2006]

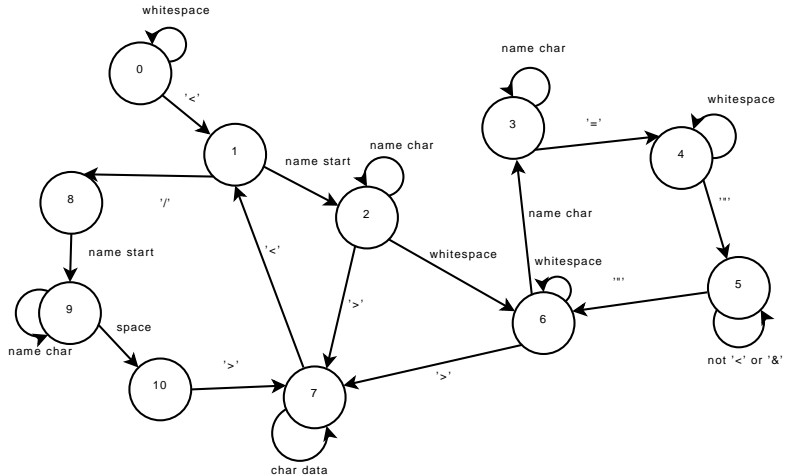
High Performance XML Processing Approaches

- Look-aside buffers/String caching [gsoap, XPP]
- Trie data structure with schema-specific parser [Chiu et al 02, Engelen 04]
- One pass table-driven recursive descent parser [Zhang et al 2006]
- Pre-scan and schedule parser [Lu et al 2006]
- Parallelized scanner, scheduled post-parser [Pan et al 2007]

Token-Scanning With a DFA

- DFA-based table-driven scanning is both popular and fast
 - (or at least performance-competitive with other techniques)
- Input is read *sequentially* from start to finish
 - Each character is used to transition over states in a DFA
 - Transition may have associated actions
 - Supports languages that are not “regular”
- Commonly used in high performance XML parsers, such as TDX (C) and Piccolo (Java)
 - Amenable to SAX parsing
 - PIXIMAL-DFA uses this approach

DFA Used in PIXIMAL-DFA



Parallel Scanning With a DFA?

- DFA-based scanning \implies sequential operation
- Desire: run multiple, concurrent DFAs throughout the input
 - Generally not possible because the start state would be unknown

Overcoming Sequentiality With an NFA

- Problem: start state is unknown
- Solution: assume every possible state is a start state
 - Construct an NFA from the DFA used in PIXIMAL-DFA
 - Such an NFA can be applied on any substring of the input
- PIXIMAL-NFA is the parser that does all of this:
 - Partition input into segments
 - Run PIXIMAL-DFA on the initial segment
 - Run NFA-based parsers on subsequent partition elements
 - Fix up transitions at partition boundaries and run queued actions

PIXIMAL-NFA's Parameters

- *split_percent*:
 - The portion of input to be dedicated to the first element of the partition, expressed as a percentage of the total input length
- *number_of_threads*:
 - The number of threads to use on a run

Preliminary Questions

- Is there enough memory bandwidth to allow multiple automata to concurrently feed each thread its input?
- Processing each character along several paths through the NFA is costly: how does this work scale with the size of the initial DFA?
- Does the overhead of queuing the NFA actions cost a reasonable amount compared with the cost of DFA-parsing the first partition element?

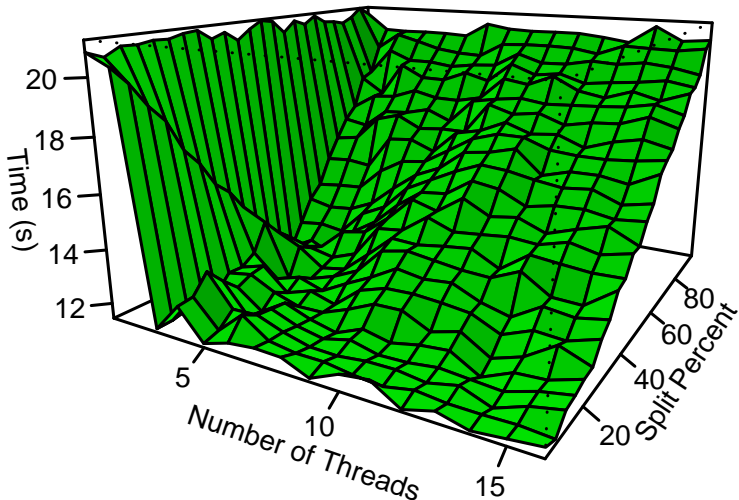
Memory Bandwidth Test

- Models the work of partitioning the input the way PIXIMAL-NFA does
 - File I/O is via `mmap (2)`
- A thread is created for each partition element which accumulates each character
- A variety of *split_percents* and *number_of_thread* are chosen
 - Total time to read a large input a fixed number of times is measured
 - Input file is `SwissProt.xml`, which is 109 MB in size

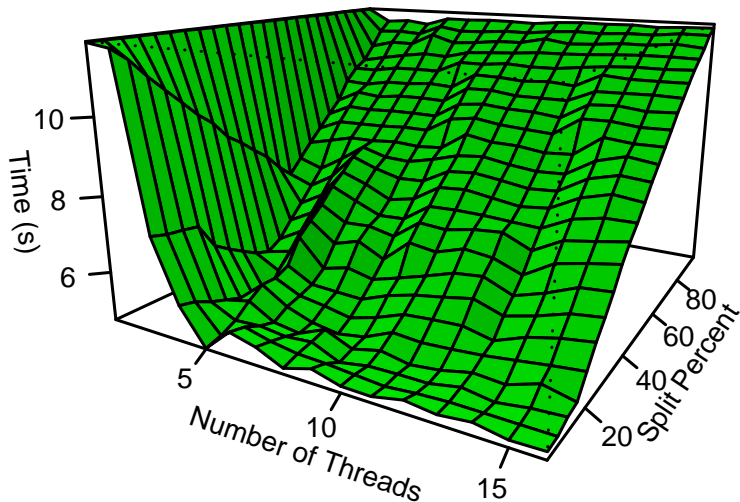
Memory Bandwidth Test – Experimental Setup

- Run several machines, each from a homogeneous class running 64-bit versions of Linux
 - 2× uniprocessor: 3.2 Ghz Intel Xeon (uniprocessor), 4 GB RAM, Linux kernel 2.6.15, GNU Lib C 2.3.6, GCC 4.0.3
 - 2× dual core: 2.66 Ghz Intel Xeon 5150 (dual core) CPUs, 8 GB RAM, Linux kernel 2.6.18, GNU Lib C 2.3.6, GCC 4.1.2
 - 2× quad core: 2.33 Ghz Intel Xeon E5354 (quad-core) CPUs, 8 GB RAM, Linux kernel 2.6.18, GNU Lib C 2.3.6, GCC 4.1.2
- 4 nodes used from the 2× UP cluster, 10 from each of the other two
- Results for each class are averaged across all runs

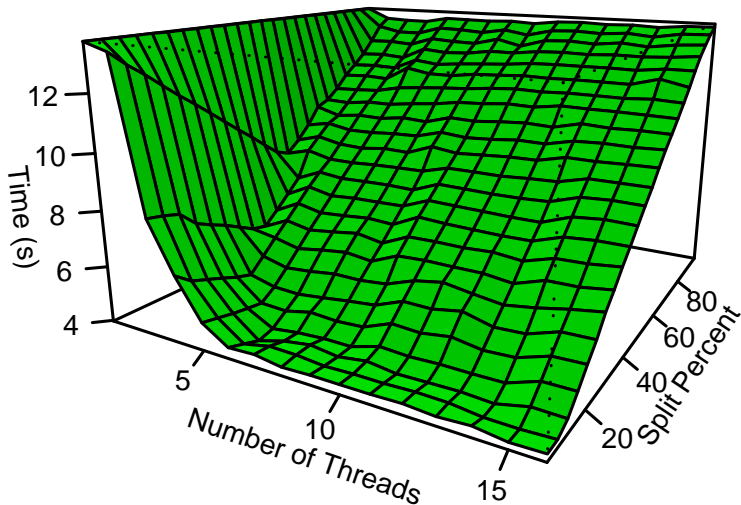
2× UP Overall Results



2× DC Overall Results



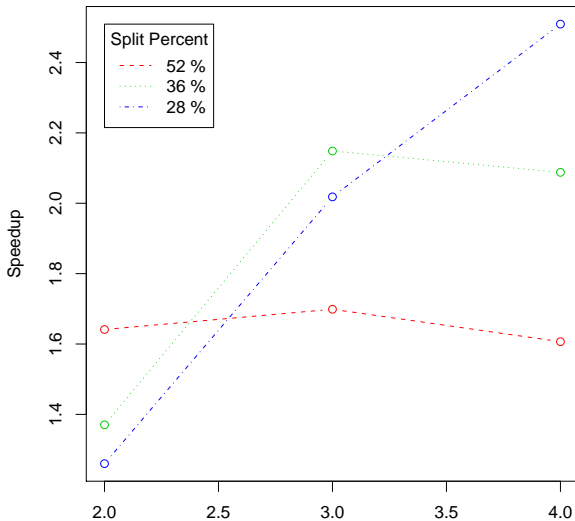
2× QC Overall Results



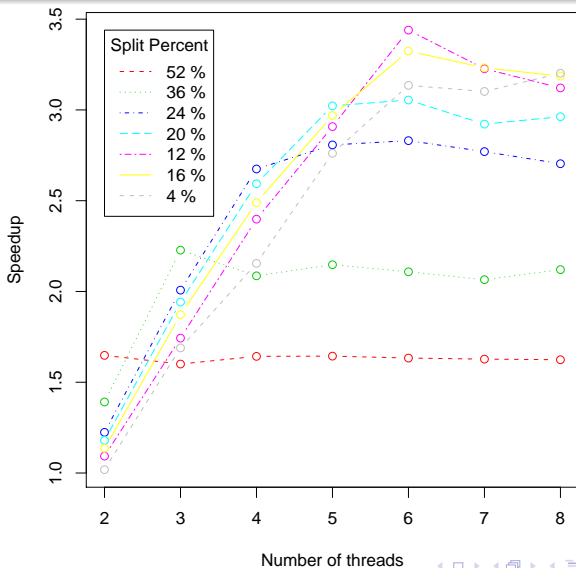
Conclusions From Overall Results

- Even when doing very little per-character processing, performance gains possible by adding threads
- Returns do diminish rapidly
- More cores lead to smoother results
- Adding “too many” threads does not hurt performance in this test
- How much gain in terms of speedup?
 - Calculated by $\frac{T_1}{T_P}$

2× DC Speedup For Best *split_percents*



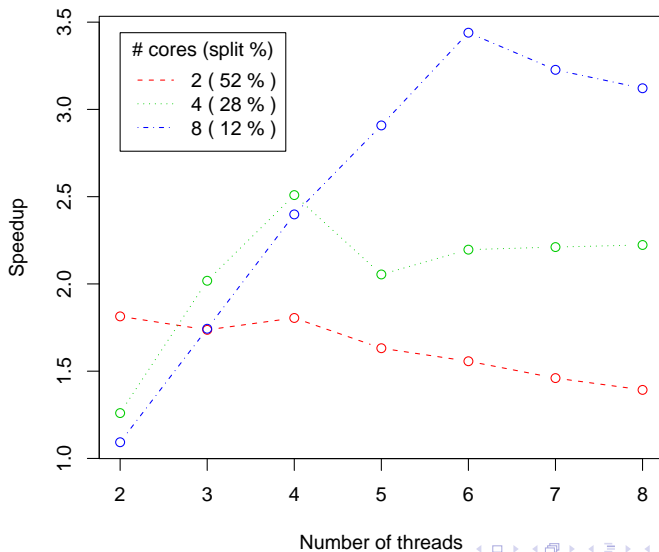
2× QC Speedup For Best *split_percents*



Conclusions From Speedup Cross Sections

- Reaffirmation that speedup is possible
- Returns diminish for these machines at around 6 threads
- Overall, access to main memory is not an immediate bottleneck
- Putting the results from the best *split_percents* for each architecture...

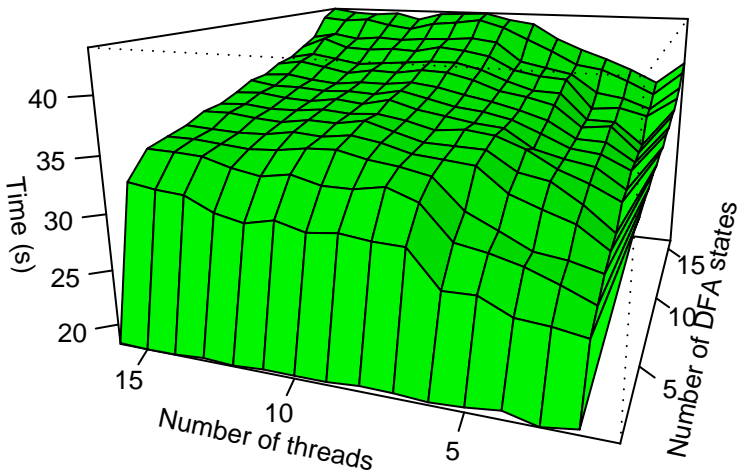
Comparison of Best *split_percent* Per Class



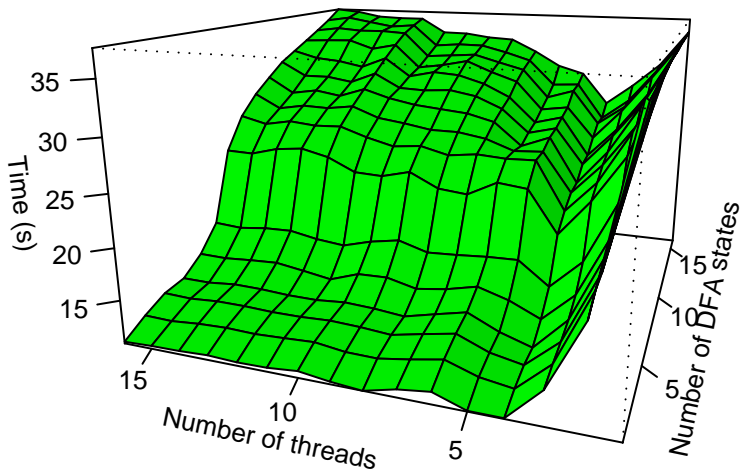
State Scalability Test

- Models the additional work done by the NFA threads by following multiple execution paths through the table
- Each NFA thread now must remember the state and calculate the next state for each character and for each start state
 - The DFA need only remember and calculate one state per input character
- Does not model the memory used, actions stored, or garbage state elimination

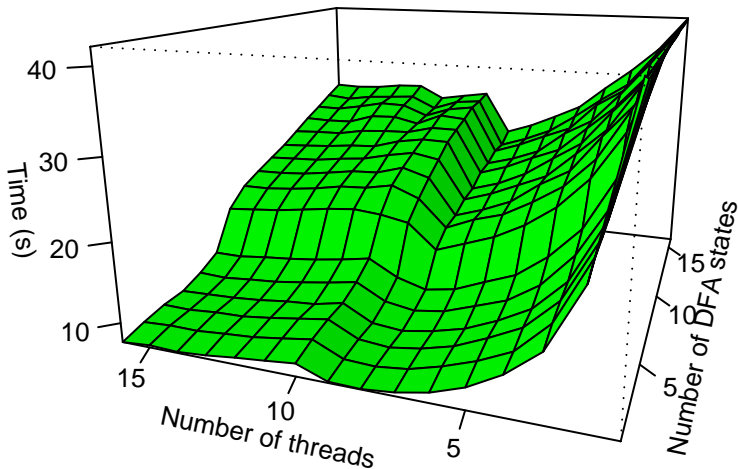
2× UP Overall Raw Results



2× DC Overall Results – Best Times



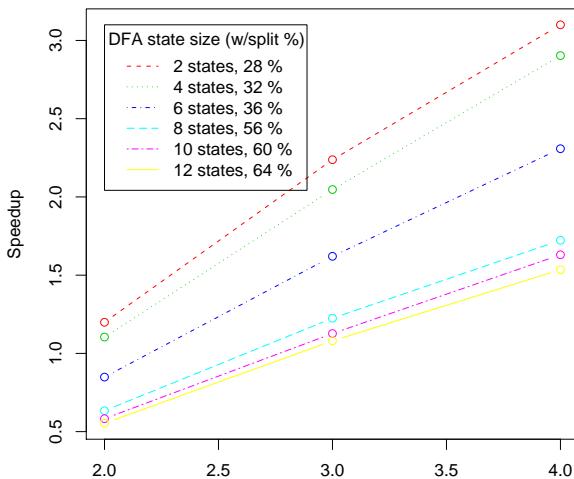
2× QC Overall Results – Best Times



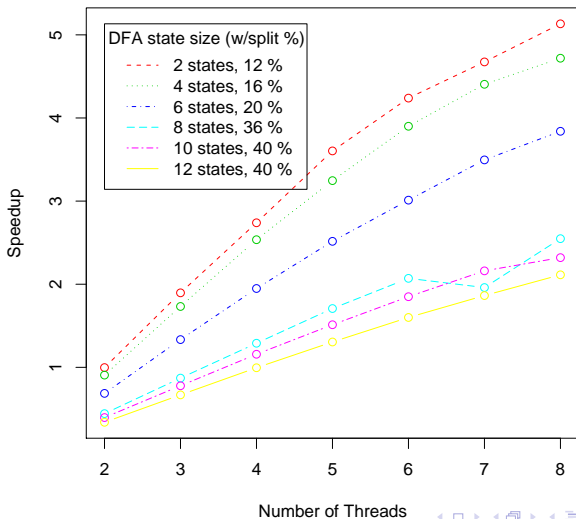
Conclusions From State Scalability Overall Results

- Two major conclusions:
 - The speedup on the $2\times$ quad-core machines appears stable as the number of threads increases
 - There is a significant steepening when the DFA has 6-7 states
- Performance reaches its max when the number of threads match the number of processing cores available
 - Each new thread adds substantial extra work compared with the memory bandwidth test
- Plotting speedup for certain *split_percents*

2× DC – Best Speedup for DFA Sizes



2× QC – Best Speedup for DFA Sizes



Conclusions From State Scalability Test

- The extra work of pushing characters through the multiple execution paths of the NFA is not in itself a limiting factor
- There is a “sweet spot” for DFA size: around 6-7 states which allows for the greatest language complexity and the best scalability
 - This is a crossover point where the $O(N)$ extra NFA work overcomes the $O(1)$ work of simply reading the input

Thank you for your time.

Questions?

Extra Slides

The following slides are additional and not part of the presentation.